

Transparent TSN for Agnostic End-hosts via P4-based Traffic Characterization at Switches

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Motivation

Challenges in Mission-Critical Systems







Background

Standard Ethernet is insufficient

Standard Ethernet

Best-effort packet delivery

No inherent latency guarantees

Optimized for general-purpose

Problem:

Standard Ethernet cannot

- Satisfy MCS timing and bandwidth constraints
- Inherently solve its packet loss
- Offer security for resource constrained devices

Time-Sensitive Networking (TSN)

Deterministic delivery of data

Traffic-shaping and QoS

Low-latency and high bandwidth

Problem:

TSN....

- Is built upon several standards
- Networks have compatibility issues
- Requires careful bandwidth allocation
- Needs to overcome security implications

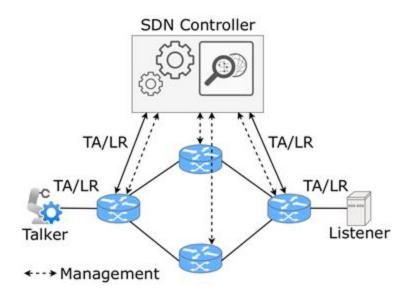
Problem statement and Research Question

End-hosts need to know **how** to schedule their transmission (according to the network's timing and rules)

= TSN-Awareness

Research Question:

How can we eliminate end-host dependencies in TSN and facilitate traffic flow management?



Centralized Control via SDN

- + Simplified resource management
- + Facilitates dynamic traffic control
- Latency concerns
- Potential bottleneck as network grows

Approach

Transparent TSN

Shift TSN responsibilities from end-hosts to network infrastructure

- Traffic classification
- Traffic identification
- Network management interaction



SDN Controller
End-hosts
Switches

- + Lower complexity
- + Lower network management overhead
- Reduced latency
- + More scalable and robust network
- + Anomaly detection

Approach

Decentralized In-Switch Monitoring

Edge switches:

- Disassemble packets
- Extract traffic characteristics
- Determine traffic type locally

$$R(k) = \frac{\sum_{t=1}^{N-k} (x_t - \bar{x})(x_{t+k} - \bar{x})}{\sum_{t=1}^{N} (x_t - \bar{x})^2} \quad \forall N, \bar{x}, k \in \mathbb{N}$$

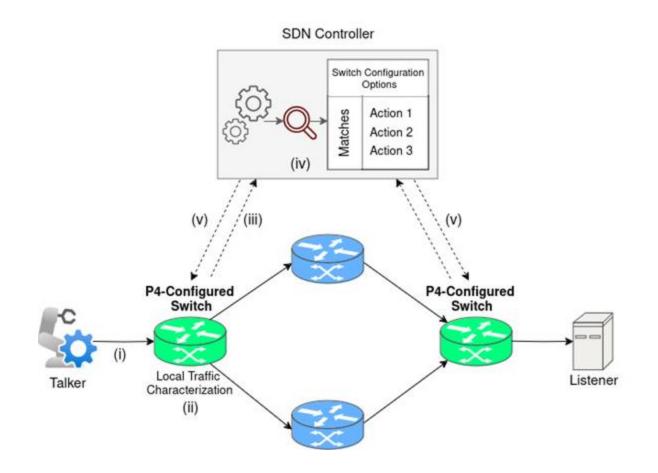
R(k): Autocorrelation coefficient at lag k.

N: Total number of frame inter-arrival times.

 $x_t \ge 0$: Inter-arrival time of the t-th frame.

 \bar{x} : Mean inter-arrival time.

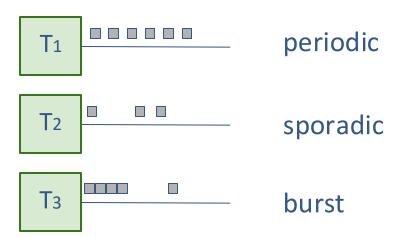
k: Lag, inter-arrival times displacement.



Implementation

System Setup

- Network setup in Mininet 2.3.1ba
- Simple Switch Behavior Model v2 (BMv2)
- Switch programming using P4
- Multiple scripts for traffic generation
- Calculate the autocorrelation to classify traffic







Implementation

Traffic Characterization

Problem No. 1:

P4 does **not maintain state** between packet processing operations.

Problem No. 2:

Which algorithm to pick for traffic characteristics when P4 can only use **limited** arithmetic operations?

Problem No. 3:

How much traffic should be monitored / stored for classification?

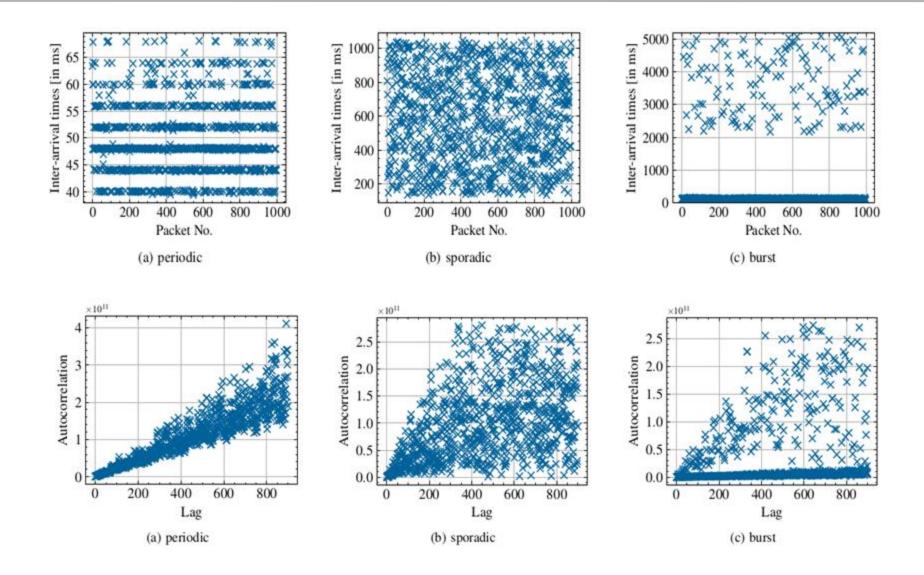
Counter Register

| Flow Iteration | Inter- Arrival | Per Flow Packet | Periodic | Sporadic | Burst | |
|------------------------|-------------------|--------------------|-----------------------------|-----------------------------|-------|-----------------------------|
| Flow Register Per Flow | | | | | | |
| Src Address | Dst Address | Prev Timestamp | Inter- Arrival Time 1 | Inter- Arrival Time 2 | | Inter- Arrival Time N |

Autocorrelation Register

| | Src Dst Address Address | Det | Auto- | Auto- | Auto- | | Auto- | ı |
|---|----------------------------|---------|-------------|-------------|-------------|--|-------------|---|
| I | | Address | correlation | correlation | correlation | | correlation | ı |
| l | | | 1 | 2 | 3 | | N-1 | ı |

Results Traffic Characterization with P4 on the Switch



Evaluation

Traffic Characterization with P4 on the Switch

Sliding Windows Evaluation:

Register Size N = 500 / 200 / 100 / 50

Precision and Recall for varying register sizes

| Register Size N | 50 | 100 | 250 | 500 |
|-----------------|------|------|------|-----|
| Precision | 0.7 | 0.67 | 0.73 | 0.8 |
| Recall | 0.74 | 0.67 | 0.73 | 1 |

Confusion Matrix

200 periodic and 200 non-periodic traffic flows

| Sent vs. Characterized | Periodic | Non-Periodic |
|------------------------|----------|--------------|
| Periodic (sent) | 200 (TP) | 0 (FN) |
| Non-Periodic (sent) | 51 (FP) | 149 (TN) |

Forwarding Latency

tshark to capture ingress and egress ports on the switch: **0.22 ms** per packet

Conclusion

Our approach:

- Reduces communication overhead / network load
- Reduces latency compared to traditional / SDN approaches
- Remove end-host dependencies for reservation requests
- In-Switch anomaly detection for traffic classes

Network participants reap the benefits of TSN

"Costs" only 0.2 ms per packet

Towards autonomous switches for a flexible network

Facilitates the switch from legacy systems to time-sensitive networks

We taught resource-constrained devices (switches) to characterize time-sensitive traffic locally!



Thank you!

Time for questions

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Abstract—Mission-critical actions currently face a transition from ligacy network protected to advanced time-mention on-working CDNs standards. This guarantees related and deter-nished communication using off-the-shelf filter not equipment. However, end-boats must be TSN-awater and may pose security risks by arthravely survailmenting resources, beingsteing central instances like a enforced-school activarising collection presents a protecting unbelow. This return is network management parents a protecting unbelow. This return is excessed to regarding latency in communication between whiches and the constraint, a welf-ary protected that renders TSN transparent to end-hosts, eliminating the need for their involvement in resource reservations. We embed packet processing legic in Pennshell TSN witheless to characteries network traffic intelligently. This enables withless to allocate action by resources and continuously and adjust real-time allocate action between traffic intelligently. This enables withless to allocate action for some contents. toal tearrier merket, vature untiligenor, the relation's victime to allocal author's resources automorphism and algust real-time traffic handling mechanism. Leveraging P4 storage structures introduces attenduces for traffic characterization computing which the inhumoity stations P4 language. Our experiments of the contract of the co

nere 9.2 ms forwarding latency per packet.

Inder Terms—P4 Programming, Traffic Characterization,
Software-Defined Networking, Time-Sensitive Networking, Real-Time Traffic Management

latency settings, face the challenge of adapting to modern system environments' increasing betweeposein and against In work to autonomously manage and adapt to traffic demands response, time-sensitive networking (TSN) has emerged, offer and communicate response, achieving transparent TSN. By ing deterministic and reliable communication in low-knowcy identifying various traffic types in real-time, the network can settings that align well with the demands of modern system dynamically allocate resources, reduce the burden on end environments. For example, network infrastructure in smart hosts, and minimize messaging frequency to the controller manufacturing facilities must meet stringent quality of service. Hence, our solution for transparent TSN involves transferring (QoS) and time-sensitive domands, where integrity in real-time. TSN tasks and network management responsibilities from the data exchange is crucial for scamics operations. However, as and-hosts and the TSN controller to the network or telesthere is a shift towards adopting ont-effective commercial off. Empowering switches to manage and identify staffic locally the shell (COTS) hardware and software over specialized field eliminates the need for end-host TSN-awareness. Further bases, orchestrating interactions among network components more, autonomous traffic type characterization at switches remains complex in real-time applications. Implementing TSN melaces messaging frequency between the TSN controller and using COTS hardware and software can ensure QoS, but mrkches, minimizing network load and reducing latency. it requires end-hosts to be TSN-aware, thus restricting the network's flexibility.

In response, ongoing sessanth explores the feasibility of central control in time-sensitive networks. Network manage ment can be facilitated, e.g., through the bolistic view of a centralized entity, similar to the concept of software-defined networking (SDN) [I]. SDN allows modification and spec-fication of network behaviour through software, effectively separating the control plane from the data plane. A central entity, known as the SDN controller, dynamically manages the traffic and configures the forwarding behaviour of the switches based on the specific application requirements. However, SDN also introduces challenges, especially in heterogeneous network environments, such in TSN. The central (TSN) controller must accommodate diverse device characteristics and network policies. This may smalt in latencies of up to 0.3 ms per packet [2] and significantly impact network operations' redtime nature. Additionally, the dynamic traffic and resource or chestration in TSN networks, involve fraguest communication between network components and the TSN controller [1]. [4] As network stuffic increases, the controller must process solumes of messages, which may become a bottleneck in the

These challenges show the segent need for a transparent TSN mechanism that guarantees dynamic, determinantic, and lowlatency communication for the desired QoS without mandating TSN-capable end-hosts. In this work, we propose to M ISSION-CRITICAL networks, characterized by their shift TSN tasks and configuration responsibilities from the determinatio and reliable communication in low-network's periphery (the end-hosts) to its internal network components to address this challenge. This enables the netCornelia Brülhart

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